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UNITED STATES PATENT AND TRADEMARK OFFICE

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BEFORE THE BOARD OF PATENT APPEALS  
AND INTERFERENCES

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*Ex parte* DUANE ARLYN AVERILL, RUSSELL DEAN HOOVER,  
DAVID ALAN SHEDIVY, and MARTHA ELLEN VOYTOVICH

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Appeal 2009-002615  
Application 10/760,431  
Technology Center 2100

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Decided: January 22, 2010

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Before JOSEPH L. DIXON, THU A. DANG, and  
CAROLYN D. THOMAS, *Administrative Patent Judges*.

DIXON, *Administrative Patent Judge*.

DECISION ON APPEAL

The Appellants appeal under 35 U.S.C. § 134(a) from a Final Rejection of claims 1-22. We have jurisdiction under 35 U.S.C. § 6(b).

We affirm-in-part.

## I. STATEMENT OF THE CASE

### *The Invention*

The invention at issue on appeal relates to a method and apparatus for selectively determining whether to send cache line invalidation messages to a particular device cache (Spec. 6-7).

### *The Illustrative Claim*

Claim 1, an illustrative claim, reads as follows:

1. A digital data processing system, comprising:
  - a memory;
  - at least one processor having at least one associated cache for temporarily caching data from said memory;
  - at least one device having a device cache, said device cache having a fixed number of slots for caching data, said fixed number being greater than one, each slot caching a cache line of data; and
  - a cache coherency mechanism, said cache coherency mechanism including a cache line state directory structure, said cache coherency mechanism selectively determining whether to send cache line invalidation messages to said at least one device using state information in said cache line state directory structure, wherein at least a portion of said cache line state directory structure associated with said at least one device contains exactly said fixed number of cache line entries, each entry having a fixed correspondence to a unique respective one

of said fixed number of slots for caching data of said device cache.

*The References*

The Examiner relies on the following references as evidence:

Baumgartner	US 6,108,764	Aug. 22, 2000
Carpenter	US 6,115,804	Sept. 5, 2000
Arimilli	US 2003/0009643 A1	Jan. 9, 2003

*The Rejections*

The following rejections are before us for review:

Claims 1, 2, and 5-8 stand rejected under 35 U.S.C. § 103(a) as being unpatentable over Baumgartner in view of Arimilli.

Claims 3, 4, and 9-22 stand rejected under 35 U.S.C. § 103(a) as being unpatentable over Baumgartner in view of Arimilli, and further in view of Carpenter.

## II. ISSUES

Have the Appellants shown that the Examiner erred in identifying that the combination of Baumgartner and Arimilli teaches and fairly suggests a cache coherency mechanism “selectively determining whether to send cache line invalidation messages to said at least one device using state information in said cache line state directory structure,” as recited in independent claim 1?

Have the Appellants shown that the Examiner erred in identifying that the combination of Baumgartner and Arimilli teaches and fairly suggests the limitations recited in claim 14?

### III. PRINCIPLES OF LAW

#### *Prima Facie Case of Unpatentability*

The allocation of burden requires that the United States Patent and Trademark Office (USPTO) produce the factual basis for its rejection of an application under 35 U.S.C. §§ 102 and 103. *In re Piasecki*, 745 F.2d 1468, 1472 (Fed. Cir. 1984) (citing *In re Warner*, 379 F.2d 1011, 1016 (CCPA 1967)). The Appellants have the burden on appeal to the Board of Patent Appeals and Interferences (BPAI) to demonstrate error in the Examiner's position. See *In re Kahn*, 441 F.3d 977, 985-86 (Fed. Cir. 2006) (citing *In re Rouffet*, 149 F.3d 1350, 1355 (Fed. Cir. 1998)).

#### *Obviousness*

The question of obviousness is resolved on the basis of underlying factual determinations including (1) the scope and content of the prior art; (2) any differences between the claimed subject matter and the prior art; and (3) the level of skill in the art. *Graham v. John Deere Co.*, 383 U.S. 1, 17 (1966).

### IV. FINDINGS OF FACT

The following findings of fact (FFs) are supported by a preponderance of the evidence.

#### *Baumgartner*

1. Baumgartner discloses a cache directory structure for storing the information relating to the cache coherency:

Coherence directory 50 stores indications of the system memory addresses of data (e.g., cache lines) checked out to caches in remote nodes for which the local processing node is

the home node. The address indication for each cache line is stored in association with an identifier of each remote processing node having a copy of the cache line and the coherency status of the cache line at each such remote processing node. Possible coherency states for entries in coherency directory 50 are summarized in Table VI.

(Col. 7, ll. 59-67.)

2. Baumgartner only mentions once “invalidation” in TABLE VI, but does not teach how to determine whether to send cache line invalidation messages using the cache line directory. (Col. 8, Table VI, ll. 27-51.)

*Arimilli*

3. Arimilli discloses a method and system of a Non-Uniform Memory Access (NUMA) computer system for improving node controller queue by implementing a reissue protocol for request of remote data (¶[0013], fig. 1). Arimilli also discloses that a structure of cache directory includes a Local Memory Directory (LMD) storing an indication regarding whether the data, resided in an associated system memory, are cached in one or more remote nodes, and a Remote Memory Directory (RMD) storing an indication regarding whether the data from the system memories in other nodes is cached in the associated RMD (¶[0037]). Arimilli further discloses the details of the cache directory 140 includes a number of sets (row) 142. Each set 142 contains a number of directory entries 144 for storing the address tag and coherency state of a corresponding cache lines (¶¶[0056] & [0057], Fig. 4).

4. Arimilli further teaches in certain circumstances such that the data is still being actively used in the remote node, or the cache line’s snoop queues are all busy, or if the target cache lines are held in the associated L2 or L1 caches, the Flush query may be ignored (¶¶[0091] & [0099]).

*Carpenter*

5. Carpenter discloses a NUMA computer system for plural cache hierarchies coupled by a local interconnection that can hold concurrently an unmodified copy of a particular cache line (col. 2, l. 56 to col. 3, l. 10).

6. Carpenter also discloses how the node controllers will operate according to the Table X (col. 12, ll.1-45, Table X).

V. ANALYSIS

“[T]he examiner bears the initial burden, on review of the prior art or on any other ground, of presenting a *prima facie* case of unpatentability. . . . If that burden is met, the burden of coming forward with evidence or argument shifts to the applicant.” *In re Oetiker*, 977 F.2d 1443, 1445 (1992).

The Examiner sets forth a detailed explanation of a *prima facie* case of unpatentability in the Examiner’s Answer. Therefore, we look to Appellant’s Briefs to show error in the proffered *prima facie* case.

*The Common Feature in Claims*

Independent claim 1, recites, *inter alia*, “cache coherency mechanism selectively determining whether to send cache line invalidation messages to said at least one device using state information in said cache line state directory structure”. Independent claims 9 and 19 contain those similar limitations recited in claim 1. The recited limitations have the same scope of selectively determining whether to send cache line invalidation messages a device by using the information in the cache line state directory.

*35 U.S.C. § 103(a) rejections*

With respect to independent claim 1, Appellants contend that neither Baumgartner nor Arimilli (nor Carpenter) discloses the feature of “selectively determining whether to send cache line invalidation messages to said at least one device” (App. Br. 12-13).

The Examiner states that Baumgartner teaches the argued limitation (Ans. 4 and Ans. 15). The Examiner also states that “Arimilli clearly discloses such limitation in at least paragraph [0056]. . . .” (Ans. 12), and “Arimilli similarly discloses the cache coherency mechanism of Appellant’s claimed invention.” (Ans. 15.)

We disagree with the Examiner’s reading of the references. We find that the Baumgartner reference relied upon by the Examiner only discloses that a cache directory stores the memory addresses of data and each cache line with an identifier relating to the remote nodes and the corresponding coherency status (FF 1). However, the Baumgartner reference does not teach how to send invalidation messages using the cache line directory, let alone selectively sending invalidation messages (FF 2).

We next look to the teachings of Arimilli. The Examiner in the Final Rejection has merely relied upon Arimilli for the limitation of the cache line state directory associated with at least one device containing a fixed number of cache line entries (Ans. 4). Furthermore, the Examiner has only provided a limited discussion of the teachings of the Arimilli reference (*id.* at 15) which is not sufficient to teach or fairly suggest the specific limitations recited in claim 1. We also note that the Arimilli reference additionally teaches that the cache flush queue may be ignored in certain circumstances (FF 4), however the Examiner does not address those teachings and how the

merits of the teachings relates to the claimed invention. Thus, we are left to speculate how the teachings may apply to the claimed limitations. We, therefore, find the Examiner's position in reliance upon the teachings of Arimilli to be untenable.

Because we agree with at least one of the Appellants' contentions, we find that the Examiner has not made a requisite showing of obviousness as required to teach or fairly suggest the invention as recited in claim 1 by the combination of Baumgartner and Arimilli. The rejection of dependent claims 2 and 5-8 contains the same deficiency. The Appellants, thus, have demonstrated error in the Examiner's *prima facie* case for obviousness of the subject matter of claims 1, 2, and 5-8.

The independent claims 9 and 19 contain the similar limitations to those found in independent claim 1. The Appellants present similar arguments as set forth with respect to independent claim 1 in response to the rejections of independent claims 9 and 19 (App. Br. 13-14).

The Examiner states that Carpenter discloses the argued limitation in the Final Rejection (Ans. 7).

We disagree with the Examiner's reading of the reference. We find that Carpenter teaches a NUMA computer system can store concurrently a cache line of data in different cache hierarchies (FF 5). However, we find Carpenter does not teach the argued limitation as the Examiner found. We find the portion relied upon by the Examiner only discloses how the node controller changes the coherence directory according to the Table X (FF 6) and does not mention sending invalidation messages.

As in our discussion with respect to independent claim 1, we similarly find that the Appellants have demonstrated error in the Examiner's *prima*

facie case for obviousness of the subject matter of independent claims 9 and 19. The rejection of dependent claims 3, 4, 10-13, and 20-22 also contains the same deficiency. Hence, the Appellants' argument persuades us that the Examiner erred in rejecting claims 1-13, and 19-22.

We, therefore, cannot sustain the rejection of claims 1-13, and 19-22 under 35 U.S.C. § 103.

With respect to claims 14-18, we note that claims 14-18 do not contain the limitation "cache coherency mechanism selectively determining whether to send cache line invalidation messages to said at least one device using state information in said cache line state directory structure," as recited in claim 1. We also note that the Appellants do not specifically argue claims 14-18, but instead argue claims 14-18 with the group of claims 3, 4, 12-18, and 21 (App. Br. 15-16). Claim 14 recites:

[W]herein said first portion of said cache line state directory structure contains a plurality of cache entries, each entry having a fixed correspondence to a respective set of real addresses; wherein said second portion of said cache line state directory structure contains exactly said fixed number of cache line entries, each entry having a fixed correspondence to a unique respective one of said fixed number of slots for caching data of said device cache.

(App. Br. 23.)

The Appellants contend that "the cache line directory structure contains multiple portions, one of which is organized having the one-to-one correspondence with cache slots discussed above, and the other of which has entries corresponding to real addresses." (App. Br. 15).

We disagree with the Appellants' contention. We find Baumgartner teaches that one cache line directory contains indications of the cache line

addresses checked out to caches in remote nodes (FF 1). We also find Arimilli teaches that a cache line directory structure includes two portions of cache line directory: LMD and RMD. The address tag and coherency state of a corresponding cache line (slot) are stored in the directory entries 144 in the LMD or RMD (FF 3). We thus find that combining the well-known element of cache line directory of Baumgartner with the well-known technique of separate cache line directory for addresses and corresponding entries of cache lines taught by Arimilli is nothing more than a “predictable use of prior art elements according to their established functions.” *KSR Int'l Co.*, 550 U.S. at 417.

Accordingly, we sustain the Examiner's obviousness rejection of the dependent claim 14. We also sustain the Examiner's obviousness rejections of dependent claims 15-18, which have analogous wording and are not separately argued and fall with their respective base claims. 37 C.F.R. § 41.37 (c)(1)(vii). *See In re Nielson*, 816 F.2d 1567, 1572 (Fed. Cir. 1987).

## VI. CONCLUSION

We conclude that the Appellants have shown that the Examiner erred by failing to identify that the combination of Baumgartner and Arimilli teaches and fairly suggests “cache coherency mechanism selectively determining whether to send cache line invalidation messages to said at least one device using state information in said cache line state directory structure,” as recited in independent claim 1.

We also conclude that the Appellants have not shown that the Examiner erred in identifying that the combination of Baumgartner and Arimilli teaches and fairly suggests the limitations recited in claim 14.

## VII. ORDER

We reverse the obviousness rejections of claims 1-13 and 19-22 under 35 U.S.C. § 103(a).

We affirm the obviousness rejections of claim 14-18 under 35 U.S.C. § 103(a).

No time period for taking any subsequent action in connection with this appeal may be extended under 37 C.F.R. § 1.136 (a). *See* 37 C.F.R. § 1.136(a)(1)(iv).

## AFFIRMED-in-PART

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